

# Local Complexity Zoo

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This is a short guide to the complexity classes and concepts in-scope for class and problem sets. Languages are sets of strings and complexity classes are sets of languages.

## 1 Inclusions & Separations: Infinitely-Often Classes (i.o. $\mathcal{C}$ )

Often, we will want to decompose a language  $\mathcal{L} \subseteq \{0, 1\}^*$  into  $n$ -bit *slices*  $\mathcal{L}_n = \mathcal{L} \cap \{0, 1\}^n$ . Given a complexity class  $\mathcal{C}$ , i.o.  $\mathcal{C}$  is the class of languages  $\mathcal{L}$  for which there is a language  $\mathcal{L}' \in \mathcal{C}$  such that  $\mathcal{L}_n = \mathcal{L}'_n$  for infinitely many input lengths  $n$ . This is a way to make inclusions *weaker* — if  $\text{SAT} \in \text{i.o. P}$  then this could be quite useless. It could be that the input lengths  $n$  for which the SAT algorithm agrees with SAT are too far apart to matter. It is also a way to make separations *stronger* — if complexity class  $\mathcal{D} \not\subseteq \text{i.o. } \mathcal{C}$  then at **every** input length past a certain point, there exists an error witnessing that some language  $H \in \mathcal{D}$  is not decidable in  $\mathcal{C}$ . For an extended discussion of these issues see the introduction of [FS17].

## 2 Problems

**Definition 1** (Minimum Circuit Size Problem [KC00]). Given as input the truth table of an  $n$ -bit Boolean function  $f$  and a number  $s \in \{1, \dots, 2^n\}$ , accept if  $f$  can be computed by a Boolean circuit with at most  $s$  gates.

## 3 Complexity Classes

**Definition 2** (Zero-Error Probabilistic Time (Ex 7.6 of [AB09])). A language  $\mathcal{L}$  is in  $\text{ZPTIME}[T(n)]$  iff there exists a  $T(n)$ -time probabilistic TM  $M$  with outputs in  $\{0, 1, ?\}$  such that, for every input, if it answers with 0 or 1 the answer is correct, and the probability of answering “I don’t know” ( $= ?$ ) is bounded:

$$\forall x \in \{0, 1\}^* \quad \Pr[M(x) \in \{\mathcal{L}(x), ?\}] = 1 \wedge \Pr[M(x) = ?] \leq 1/2$$

## References

- [AB09] Sanjeev Arora and Boaz Barak. *Computational Complexity - A Modern Approach*. Cambridge University Press, 2009. ISBN: 978-0-521-42426-4. URL: <http://www.cambridge.org/catalogue/catalogue.asp?isbn=9780521424264>.
- [FS17] Lance Fortnow and Rahul Santhanam. “Robust simulations and significant separations”. In: *Inf. Comput.* 256 (2017), pp. 149–159. DOI: 10.1016/j.ic.2017.07.002. URL: <https://doi.org/10.1016/j.ic.2017.07.002>.
- [KC00] Valentine Kabanets and Jin-yi Cai. “Circuit minimization problem”. In: *Proceedings of the Thirty-Second Annual ACM Symposium on Theory of Computing, May 21-23, 2000, Portland, OR, USA*. Ed. by F. Frances Yao and Eugene M. Luks. ACM, 2000, pp. 73–79. DOI: 10.1145/335305.335314. URL: <https://doi.org/10.1145/335305.335314>.